# Riemann-Roch on graphs, debt games and applications.

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This is a brief survey of two papers.<sup>12</sup> We cover the nature of divisors on graphs and show part of the proof of a graph-theoretic counterpart of the classic Riemann-Roch theorem. Furthermore, the theory is applied to give a new volumetric version of Kirchoff's theorem.

### 1 DOLLAR DEBT GAME

Consider any connected graph, G, with an integer weight at each vertex,  $v \in V(G)$ . Think of the weight as the amount of dollars (or debt) the vertex has. For any vertex, v, we allow two legal moves: Either v takes a dollar from each neighbor, w (i.e. v, w connected by an edge  $e = vw \in E(G)$ ), trough each edge. Or v gives a dollar to each neighbor, through each edge. The goal of the game is then to get every vertex out of debt. The natural question is of course; for which configurations is the game winnable?



Figure 1: Example of legal move.

First we need some notation. We denote the genus of any graph, G, by g := |E(G)| - |V(G)| + 1 and define the degree, deg(D), of a vertex-weighted graph, D, as the sum of all the weights (i.e. the total number of dollars). Then we have the following:

**Theorem 1.1.** If  $\deg(D) \ge g$ , then the game is always winnable. Furthermore, if  $\deg(D) \le g-1$ , then there exists some configuration for which the game is *not* winnable.

We also define r(D) to be -1 if the game is not winnable and to be a non-negative integer, k, denoting the largest number of dollars we can remove from the game (in any way), so that it remains winnable.

Finally, we define the canonical configuration *K* to be given by  $K(v) = \deg(v) - 2$ , where  $\deg(v)$ , with  $v \in V(G)$ , is the number of edges incident to *v*.

We can now express the Riemann-Roch theorem for graphs:

**Theorem 1.2.** (Riemann-Roch for graphs) For any configuration *D* on any graph *G* we have

$$r(D) - r(K - D) = \deg(D) + 1 - g$$

An immediate corollary of this is that if deg(D) = g - 1, then r(D) = r(K - D), and we see that *D* is winnable if and only if K - D is winnable.

It is easy to see how Theorem 1.1. follows from Theorem 1.2., but we will see later that these theorems can be expressed in a more interesting context!

#### 2 RIEMANN-ROCH FOR GRAPHS

Let *G* be any graph without loop edges, and let *Q* denote the Laplacian matrix. We let the *divisor group*, Div(G), be the free abelian group on V(G), and write an element  $D \in Div(G)$  as the formal sum  $\sum_{v \in V(G)} a_v(v)$ , and denote  $a_v$  by D(v). Think of *D* as a configuration on *G* from the debt game. Div(G) is the graph analogy to the divisors on a Riemann surface, but we won't go more in depth with the analogies with Riemann surfaces.

Div(*G*) has a partial ordering given by  $D \ge D'$  if and only if  $D(v) \ge D'(v)$  for all  $v \in V(G)$ . We call a divisor, *E*, effective if  $E \ge 0$  and denote by Div<sub>+</sub>(*G*) the set of effective divisors.

Furthermore, we define the degree function deg :  $\text{Div}(G) \to \mathbb{Z}$  by  $\text{deg}(D) = \sum_{v \in V(G)} D(v)$ , i.e. the sum of coefficients of *D*.

Let  $\mathcal{M}(G) = \text{Hom}(V(G), \mathbb{Z})$ , the abelian group of integer-valued functions. The Laplacian operator  $\Delta : \mathcal{M}(G) \to \text{Div}(G)$  is given by

$$\Delta(f) = \sum_{v \in V(G)} \Delta_v(f)(v)$$
$$\Delta_v(f) = \sum_{e \in E_v} (f(v) - f(w))$$

In fact, one can easily see then that  $[\Delta(f)] = Q[f]$ .

We define  $\operatorname{Div}^k(G) = \{D \in \operatorname{Div}(G) \mid \deg(D) = k\}$  as well as  $\operatorname{Div}^k_+(G) := \operatorname{Div}_+(G) \cap \operatorname{Div}^k(G).$ 

We can now define the principal divisors  $Prin(G) := \Delta(\mathcal{M}(G))$ and note that this must specially be a subgroup of  $Div^0(G)$ . And so, we can now finally define the quite interesting quotient groups, the *Picard groups* of G:

$$\operatorname{Pic}^{n}(G) = \operatorname{Div}^{n}(G)/\operatorname{Prin}(G).$$

and write [D] for the class of in  $\operatorname{Pic}^{n}(G)$  of  $D \in \operatorname{Div}^{g}(G)$ .

The Picard groups, have many interesting properties. For starters, the order  $Pic^{0}(G)$  is the number of spanning trees in G.

We now fix a base vertex  $v_0$  and define the Abel-Jacobi map

$$S_{v_0}: G \to \operatorname{Pic}^0(G), \quad S_{v_0}(v) = [(v) - (v_0)]$$

and for  $k \ge 0$  a map  $S_{v_0}^{(k)}$  :  $\operatorname{Div}_+^k(G) \to \operatorname{Pic}^0(G)$  given by

$$S_{v_0}^{(k)}((v_1) + \dots + (v_k)) = S_{v_0}(v_1) + S_{v_0}(v_2) + \dots + S_{v_0}(v_k).$$

then we have

**Theorem 2.1.** The map,  $S^{(k)}$ , is surjective if and only if  $k \ge g$ .

We define a linear equivalence relation  $D \sim D'$  if  $D - D' \in Prin(G)$ . (In terms of dollar-debt game  $D \sim D'$  if and only if there is a series of moves that connect the two). Note that it follows that D and D' have same degree and that  $Pic^n$  is the set of equivalence classes of degree n divisors on G.

One can also show that  $S^{(k)}$  is surjective if and only if every divisor of degree k is linearly equivalent to an effective divisor. From this we see that Theorem 2.1 is equivalent to Theorem 1.1!

<sup>&</sup>lt;sup>1</sup>arXiv:math/0608360v3 [math.CO] 9 Jul 2007

<sup>&</sup>lt;sup>2</sup>arXiv:1304.4259v2 [math.CO] 24 Sep 2014

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Let

$$|D| := \{E \in \operatorname{Div}(G) \mid E \ge 0, E \sim D\}$$

which we use to define the dimension as r(D) = -1 if  $|D| = \emptyset$  and  $r(D) = \max\{ s \mid |D - E| \neq \emptyset$  for all *E* of degree *s*}. Note r(D) is invariant under the equivalence, ~. (In terms of dollar-debt game r(D) is the maximum number of dollars you can remove (in any way) from the board such that it remains winnable).

Finally, the canonical divisor is given by

$$K = \sum_{v \in V(G)} (\deg(v) - 2)(v).$$

Note deg(K) = 2|E(G)| - 2|V(G)| = 2g - 2, since we get every edge twice.

Thus we can restate Theorem 1.2. in this more rigorous framework

**Theorem 2.2.** (Riemann-Roch for graphs) Let G a graph, D a divisor on G. Then

$$r(D) - r(K - D) = \deg(D) + 1 - g$$

The properties of Riemann-Roch for Graphs has a lot of overlap with those of Riemann-Roch for Riemann surfaces, however it's not one to one and one must be careful.

An easy consequence of this is Clifford's Theorem for Graphs:

**Corollary 2.3.** Let *D* be an effective divisor such that  $|K - D| \neq \emptyset$  (called special) on *G*. Then

$$r(D) \le \frac{1}{2} \deg(D)$$

#### 3 THE PROOF

There was nothing special about the set, V(G), on which defined Div(G), so we could've used any set X instead of V(G). Likewise the effective divisors and  $\text{Div}^d_+$  did not depend on anything other then the set structure, and can be generalized to any set, X. Further, if we can define  $\sim$  on Div(X) satisfying:

(E1) If  $D \sim D'$  then deg(D) = deg(G').

(E2) If  $D_1 \sim D'_1$  and  $D_2 \sim D'_2$ , then  $D_1 + D_2 \sim D'_1 + D'_2$ .

Then we can define  $|D| := \{E \in \text{Div}(X) \mid E \ge 0, E \sim D\}$ , and  $r : \text{Div}(X) \rightarrow \{-1, 0, 1, 2, ...\}$  in the same way we did for graphs.

Finally let  $\mathcal{N} = \{D \in \text{Div}(X) \mid \text{deg}(D) = g - 1, |D| = \emptyset\}$ , and *K* some divisor with degree 2g - 2.

Then we have the following generalization of Riemann-Roch, which we will take for granted:

Theorem 3.1. The Riemann-Roch equality,

$$r(D) - r(K - D) = \deg(D) + 1 - q$$

holds for all  $D \in Div(G)$  iff the following two properties hold:

(RR1) For every  $D \in \text{Div}(X)$ , there is  $\nu \in \mathcal{N}$  such that either  $|D| = \emptyset$  or  $|\nu - D| = \emptyset$ , but never both.

(RR2) For every  $D \in \text{Div}(X)$  with deg(D) = g - 1, either both  $|D| = \emptyset$  and  $|K - D| = \emptyset$ , or both are non-empty.

To prove this one needs  $v_0$ -*reduced* divisors which are divisors which satisfy  $D(v) \ge 0$  for all  $v \ne v_0$ , and that for every non-empty

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subset  $A \subseteq V(G) - \{v_0\}$  there is  $v \in A$  such that  $D(v) < \text{outdeg}_A(v)$ , where  $\text{outdeg}_A$  is the number of edges from v which end not in A.

In terms of dollar-debt game:  $v_0$  is the only vertex which can be in debt and if all  $v \in A$  were to make a lending move, some vertex of A would go into debt, for all  $A \subset V(G) - \{v_0\}$ .

**Proposition 3.2.** Fix  $v_0$  then for every divisor *D* there is a unique  $v_0$ -reduced divisor *D'* such that  $D \sim D'$ .

Given some total order  $<_P$  on V(G) we define a specefic divisor:

$$v_P = \sum_{v \in V(G)} (|\{e = vw \in E(G) \mid w <_P v\}| - 1)(v)$$

And notice that  $\deg(v_P) = |E(G)| - |V(G)| = g - 1$ , since we will see every edge and we subtract 1 at each vertex.

It turns out that  $v_P \in N$  for all total orders,  $<_P$ . Which will help us prove the following:

**Theorem 3.3.** For all divisors *D* exactly one of the following hold: (N1)  $r(D) \ge 0$ 

(N2)  $r(v_P - D) \ge 0$  for some  $<_P$ .

PROOF. Fix  $v_0$ . We may assume D is  $v_0$ -reduced by prop. 3.2. We define an order  $v_1, v_2, ..., v_{|V(G)|-1}$  (i.e.  $v_i <_P v_j$  iff i < j) iductively: If  $v_0, ..., v_{k-1}$  defined then let  $A_k = V(G) - \{v_0, ..., v_{k-1}\}$  and choose  $v_k$  so that  $D(v_k) < \text{outdeg}_{A_k}(v_k)$ .

Now for every  $v_k \neq v_0$  we have by definition of  $v_P$ 

$$\begin{split} D(v_k) &\leq \mathrm{outdeg}_{A_k} - 1 \\ &= |\{e = v_k v_j \mid v_j < v_k\}| - 1 \\ &= v_P(v_k). \end{split}$$

If  $D(v_0) \ge 0$  then  $D \ge 0$  (since it's  $v_0$ -reduced) and (N1) holds. And if  $D(v_0) \le -1$  then  $D \le v_P$ , so  $v_P - D \ge 0$  and (N2) holds. If both  $r(D) \ge 0$  and  $r(v_P - D) \ge 0$  then  $r(v_P) = r(D + v_P - D) \ge$  $r(D) + r(v_P - D) \ge 0$ , contradicting that  $v_P \in N$ .

**Corollary 3.4.** For all divisors, *D*, of degree g - 1, we have that  $D \in N$  if and only if there exists  $<_P$  on V(G) such that  $D \sim v_P$ .

PROOF. If  $v_P - D \sim E$  with  $E \ge 0$  then  $\deg(E) = \deg(v_P - D) = 0$ and so E = 0 and  $D \sim v_P$ .

Finally we can prove Theorem 2.2:

So vo

PROOF. (of Theorem 2.2) We need to show RR1 and RR2 hold. Assume  $D \in \text{Div}(G)$  with  $r(D) \ge 0$ . For all  $v \in N$  we have r(v - D) = -1 and so  $|D| \neq \emptyset$  and  $|v - D| = \emptyset$ . So RR1 holds.

On the other hand if r(D) = -1, then by [3.3], have  $r(v_P - D) \ge 0$ for some  $<_P$ , and then  $|D| = \emptyset$  and  $|v_P - D| \ne \emptyset$ . Since  $v_P \in N$ , so RR1 holds.

For RR2 it suffice to show that for all  $D \in N$  we have  $K - D \in N$ . By [3.4] we have  $D \sim v_P$  for some order  $<_P$ . Define  $<_Q$  by  $v <_Q$  $w \Leftrightarrow w <_P v$ , i.e. the reverse of *P*. Then for every *v* we have

$$v_P(v) + v_Q(v) = |\{e = vw \mid w <_P v\}| - 1$$
  
+ |{e = vw | w <\_Q v}| - 1  
= deg(v) - 2 = K(v).  
= K - v\_P and so K - D ~ K - v\_P = v\_Q \in N

#### 4 APPLICATIONS

This theory and the Riemann-Roch for graphs have many applications in various fields. One example is that we can give a profound proof of the classic Kirchhoff's Theorem:

**Theorem 4.1** (Kirchoff's Theorem). The number of spanning trees of a graph *G* is equal to *any* cofactor of the Laplacian of *G*.

First we must define *break divisors*: A break divisor, *D*, on a graph, *G*, is an effective divisor of degree g(G) such that *D* restricted to any connected subgraph *H* of *G* has the property that  $\deg(D|_H) \ge g(H)$ . It turns out break divisors have a special connection with spanning trees and the picard groups:

**Theorem 4.2** Let g be the genus of G. Then every degree g divisor is equivalent to a unique break divisor. Thus the set of break divisors on G is canonically in bijection with  $\text{Pic}^{g}(G)$ 

Since  $|\operatorname{Pic}^{g}(G)| = |\operatorname{Pic}^{0}(G)|$  and the size of  $\operatorname{Pic}^{0}(G)$  is exactly the number of spanning trees, we see that the number of break divisors on *G* equals the number of spanning trees of *G*!

Finally, before stating the theorem, we need to talk about tropical curves:

A tropical curve (or metric graph),  $\Gamma$  can be obtained from a graph *G* by assigning an edge-length  $\ell(e) \in \mathbb{R}$  to each edge  $e \in E(G)$ , and identifying *e* with the obvious line segment of that length.

Divisors on tropical curves are of the form  $\sum_{p \in \Gamma} a_p(p)$ , with only finitely many  $a_p \in \mathbb{Z}$  non-zero and p allowed to be anywhere along any edge.

Let  $f: \Gamma \to \mathbb{R}$  be any tropical rational function: a piecewise linear function with only finitely many pieces, each having integer slope. A principal divisor,  $\operatorname{div}(f)$  is then given as  $\operatorname{div}(f) = \sum_{p \in \Gamma} \operatorname{ord}_p(f)(p)$ , with  $\operatorname{ord}_p(f)$  being minus the sum of the outgoing slopes of f emanating from p.

We define  $\operatorname{Pic}^{n}(\Gamma)$  in exactly the same way as for graphs:

$$\operatorname{Pic}^{n}(\Gamma) = \operatorname{Div}^{n}(\Gamma) / \operatorname{Prin}(\Gamma)$$

Now  $\operatorname{Pic}^0$  is no longer finite group but rather a real *g*-dimensional torus.

Similarly to graphs, we can also define break divisors on a tropical curve,  $\Gamma$ . A break divisor on  $\Gamma$  is an effective divisor, D, of degree g such that D restricted to any closed connected subgraph  $\Gamma'$  has degree at least that of the genus of  $\Gamma'$ . Once again it can be shown that there is a bijection between the break divisors of  $\Gamma$  and  $\operatorname{Pic}^g(\Gamma)$ .

Furthermore, one has that *D* is a break divisor on  $\Gamma$  if and only if there exists a spanning tree *T* of *G* and an enumeration  $e_1^\circ, ..., e_g^\circ$  of  $\Gamma \setminus T$  such that  $D = (p_1) + \cdots + (p_g)$  with each  $p_i \in e_i$  (where  $e_i^\circ$  are the open edge of  $e_i$  (i.e. endpoints removed)). For a tree *T*, let  $B_T$  be the set of all divisors,  $(p_1) + \cdots + (p_g)$ , defined as above. Finally let  $C_T \subset \operatorname{Pic}^g(T)$  be the image of  $B_T$  under the map  $D \mapsto [D]$ , sending *D* to it's linear equivalence class. Then we have the following:

**Theorem 4.3** We have that  $\operatorname{Pic}^{g}(\Gamma) = \bigcup_{T \in \mathcal{T}} C_{T}$ , where  $\mathcal{T}$  is the set of all spanning trees of *G*. Furthermore, each  $C_{T} \subset \operatorname{Pic}^{g}(\Gamma)$  is a parallelotope with their relative interior disjoint.

What's more there is a natural metric on  $\operatorname{Pic}^{g}(\Gamma)$  for which  $\operatorname{vol}(C_{T}) = \prod_{e \notin T} \ell(e)$  and the volume of  $\operatorname{Pic}^{G}(\Gamma)$  is naturally related to the determinant of the Laplacian of *G*, from which one can recover Kirchoff's theorem!

#### 5 QUICK EXAMPLE

Let  $\Gamma$  be the metric graph consisting of 2 vertices connected by 3 edges of length, 2, 1 and 2. We can fix a model, *G*, for  $\Gamma$  in which all edges have length 1:



Fig. 1. Model for  $\Gamma$ 

This has the following spanning trees:



Fig. 2. Spanning trees of G

And the cell decomposition looks as follows:



Fig. 3. Cell decomposition

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